Using Adaptive Formalisms to Describe Context-Dependencies in Natural Language

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Abstract. This text sketches a method based on adaptive technology for representing context-dependencies in NL processing. Based on a previous work [4] dedicated to syntactical ambiguities and non-determinisms in NL handling we extend it to consider context-dependencies not previously addressed. Although based on the powerful adaptive formalism [3], our method relies on adaptive structured pushdown automata [1] and grammars [2] – resulting simplicity, low-cost and efficiency.

1 Introduction

Since low-complexity language formalisms are too weak to handle NL, stronger formalisms are required, most of them resource demanding, hard to use or unpractical. Structured pushdown automata are excellent to represent regular and context-free aspects of NLs by allowing them to be split into a regular layer (implemented as finite-state machines) and a context-free one (represented by a pushdown store). Such device accepts deterministic context-free languages in linear time, and is suitable as an underlying mechanism for adaptive automata, allowing handling – without loss of simplicity and efficiency – languages more complex than context-free ones.

Classical grammars may describe non-trivial interdependencies between and inside sentences: attribute-, two-level-, evolving- and adaptive- grammars. Here, context dependency is handled with adaptive grammars (which may be converted [2] into structured pushdown adaptive automata [3]) by executing adaptive actions attached to the rule being used (stating self-modifications – rule addition and deletion – to be imposed). Upon a context-dependency is detected, one of such rules is applied, and the attached adaptive action learns to handle of the context dependency by adequately changing the underlying grammar. Starting from an initial grammar, the adaptive device follows its rules until some new context-dependency is handled. Thereafter, its operation follows the modified underlying grammar until either the sentence is fully derived or no matching rule is found. Complex languages, e.g. NLs, may be handled in this way, since adaptive grammars have type-0 power [1], [2]. By converting them into adaptive structured pushdown automata, simplicity and efficiency are achieved through piecewise-regular handling of the language, validating adaptive devices as practical and efficient for NL handling [5].

$\mathbf{2}$ Illustrating Example

This example illustrates nominal agreement in Portuguese using an adaptive grammar [2] and considers: attractive agreement for adjectives placed before nouns coordinated with preposition "e" or comma (e.g. As antigas mansões e parques) and grammatical agreement for adjectives placed after such nouns (e.g. Os parques e mansões restaurados). Our adaptive grammar is defined as a 3tuple (G^0, T, R^0) where:

T = finite set of adaptive functions	P_D^0 rules used in CD situations
$G^0 = (V_N^0, V_T^0, V_C, P_L^0, P_D^0, S)$ initial grammar	$R^{\overline{0}} = P_L^0 \cup P_D^0$
$V_N^0 =$ non-empty finite set of non-terminals	The example refers to $G = (G^0, T, R^0)$:
V_T = non-empty finite set of terminals	$G^{0} = (V_{N}^{0}, V_{T}^{0}, V_{C}, P_{L}^{0}, P_{D}^{0}, S)$
$V_N^0 \cap V_T = \emptyset$	$V_N^0 = \{C, C_1, C_2, C_3, C_4, C_5, C_6, C_7, D, $
V_C = finite set of context symbols	$A, S, C_{8a}, C_{8l}, ESM, ESF, EPM, EPF$
$V^0 = V_N^0 \cup V_T \cup V_C$	$V_C = \{sm, sf, pm, pf\}$
V_N^0, V_T, V_C are disjoint sets	$V_T = \{as, e, antigas, mansões,$
$S \in V_N^0$ start symbol of the grammar	$parques, restaurados, praças, "," \}$
P_L^0 rules used in CF situations	

Context symbols sm, sf, pm, pf denote attributes: singular/plural masc./fem. D, A, S denote determinants, adjectives and nouns, respectively. Starting symbol is C. Adaptive functions dynamically handle optional elements: further nouns, a determinant, an adjective placed before/after the noun. A context-dependency is handled by an adaptive function when the noun is processed. It checks its agreement with the previous determinant and adjective. Another adaptive function enforces agreement between the adjective and multiple nouns. P_L^0 and P_D^0 are:

 $C \rightarrow \{A_1(C, C_1)\}DC_1$ $sfC_1 \rightarrow \{A_{1c}(C_1, C, ES, EF)\}ESF$ $pfC_1 \rightarrow \{A_{1c}(C_1, C, EP, EF)\}EPF$ $smC_1 \rightarrow \{A_{1c}(C_1, C, ES, EM)\}ESM$ $pmC_1 \rightarrow \{A_{1c}(C_1, C, EP, EM)\}EPM$ $C \rightarrow \{A_2(C, C_2, C_1)\}AC_2$ $sfC_2 \rightarrow \{A_{1c}(C_2, C, ES, EF)\}ESF$ $pfC_2 \rightarrow \{A_{1c}(C_2, C, EP, EF)\}ESF$ $smC_2 \rightarrow \{A_{1c}(C_2, C, ES, EM)\}ESM$ $pmC_2 \rightarrow \{A_{1c}(C_2, C, EP, EM)\}EPM$ $C \rightarrow SC_3$ $sfC_3 \rightarrow \{A_{3c}(C_6, C_1, C_2, ES, EF)\}ESF$ $smC_3 \rightarrow \{A_{3c}(C_6, C_1, C_2, ES, EM)\}ESM C_3 \rightarrow \{A_5(C_3, C_6), C_7, C_{8a}, C_{8l})\}AC_7$ $ESM \to C \quad ESF \to C$ $EPM \to C$ $EPF \rightarrow C$

 $pfC_3 \rightarrow \{A_{3c}(C_6, C_1, C_2, EP, EF)\}EPF$ $pmC_3 \rightarrow \{A_{3c}(C_6, C_1, C_2, EP, EM)\}EPM$ $C_7 \to \emptyset$ $C_3 \to \varepsilon$ $C_3 \rightarrow C_6$ $C_3 \rightarrow "e"C_4$ $C_3 \rightarrow ", "C_4$ $C_4 \rightarrow SC_5$ $smC_5 \rightarrow \{A_4(C_6, ES, EM)\}ESM$ $sfC_5 \rightarrow \{A_4(C_6, ES, EF)\}ESF$ $pmC_5 \rightarrow \{A_4(C_6, EP, EM)\}EPM$ $pfC_5 \rightarrow \{A_4(C_6, EP, EF)\}EPF$ $ESM \rightarrow C_3$ $ESF \rightarrow C_3$ $EPF \rightarrow C_3$ $EPM \rightarrow C_3$

Adaptive Actions: $A_1(XC, XC1) =$ /*remove extra determinant*/ $\{-[XC \to \{A_1(XC, XC_1)\}DXC1]\}$ $A_{1C}(x1,x2,xn,xg) =$ /*Delete transitions with improper context symbols*/ $\{-[smx1 \rightarrow \{A_{1c}(x1, x2, xn, xg)\}ESM]$ $-[sfx1 \rightarrow \{A_{1c}(x1, x2, xn, xg)\}ESF]$ $-[pmx1 \rightarrow \{A_{1c}(x1, x2, xn, xg)\}EPM]$ $-[pfx1 \rightarrow \{A_{1c}(x1, x2, xn, xg)\}EPF]$ /*ATK: dummy adaptive action. It memorizes determinant attributes*/ $+[x1 \rightarrow \{ATK(xn, xg)\}x2]\}$ $\mathbf{A}_2(\mathbf{XC}, \, \mathbf{XC2}, \, \mathbf{XC1}) =$ $\{-[XC \to \{A_1(XC, XC1)\}DXC1]$ /*disable further det. or adj.*/ $-[XC \rightarrow \{A_2(XC, XC2, XC1)\}AXC2]\} \{xn1, xn2, xm, xf:$ $\mathbf{A}_{3c}(\mathbf{xc6}, \mathbf{xc1}, \mathbf{xc2}, \mathbf{xc}, \mathbf{xn}, \mathbf{xg}) =$ $\{dn, dg, an, ag:$ /*memorize noun attributes*/ $A_4(xc6, xn, xg)$ /*fix inflexion of determinant and adjective before noun*/ $-[xc1 \rightarrow \{ATK(dn, dg)\}xc]$ $-[xc2 \rightarrow \{ATK(an, ag)\}xc]$ $+[xc1 \rightarrow \{ATK(xn, xg)\}xc]$ $+[xc2 \rightarrow \{ATK(xn, xg)\}xc]$ $A_4(xc6,xn,xg) = \{x, s*:$ /*ATKS: dummy adaptive action. It memorizes noun attributes*/ $-[x \rightarrow xc6]$ $+[x \rightarrow \{ATKS(xn,xg)\}s]$ $+[s \rightarrow xc6]\}$ $A_5(xc3,xc6,xc7,xc8a,xc8l) =$ $\{x, xsm, xsf, xpm, xpf, xn1, xn2,$ xn3, xn4, xg1, xg2, xf1, xf2, x1, x2,xaux1*, xaux2*:/* imposes attractive agreement*/ $?[x \rightarrow xc6]$ $?[xsm \rightarrow {ATKS(ES, EM)}x]$ $?[xsf \rightarrow \{ATKS(ES, EF)\}x]$ $?[xpf \rightarrow \{ATKS(EP, EF)\}x]$ $?[xpm \rightarrow \{ATKS(EP, EM)\}x]$

+[$smxc7 \rightarrow \{AT(xsm, xc7, xaux1)\}xaux1$] $+[sfxc7 \rightarrow \{AT(xsf, xc7, xaux1)\}xaux1]$ $+[pmxc7 \rightarrow \{AT(xpm, xc7, xaux1)\}xaux1]$ $+[pfxc7 \rightarrow \{AT(xpf, xc7, xaux1)\}xaux1]$ /*initializes logical agreement */ $?[xc3 \rightarrow {ATKS(xn1, EF)}xf1]$ $?[xc3 \rightarrow \{ATKS(xn2, EM)\}xm1]$ $?[xf1 \rightarrow \{ATKS(xn3, xg1)\}x1]$ $?[xm1 \rightarrow {ATKS(xn4, xg2)}x2]$ $+[pfxc7 \rightarrow$ $\{Z(xf1, x1, xaux2, xc8l)\}xaux2\}$ $+[pmxc7 \rightarrow$ $\{Z(xm1, x2, xaux2, xc8l)\}xaux2\}$ CL(xf1, xaux2, xc8l, xc7)CL(xf1,xaux2,xc8l,xc7) =/* completes logical agreement*/ $?[xf1 \rightarrow \{ATKS(xn1, EM)\}xm]$ $?[xf1 \rightarrow \{ATK(xn2, EF)\}xf]$ $+[pmxc7 \rightarrow$ $\{Z(xm, xm, xaux2, xc8l)\}xaux2\}$ EliminaPF(xm, xc7, xaux2)CL(xf, xaux2, xc8l, xc7)EliminaPF(xm,xc7,xaux2) =/* removes the pl. fem. agreement*/ ${x, y, z:$ $-[pfxc7 \rightarrow \{Z(x, y, xaux2, z)\}xaux2]\}$ Z(x,y,z,xc8l) =/*inserts a transition to a final state*/ $\{+[z \to xc8l]\}$ $AT(x,xc7,xaux1) = \{xsmx, xsfx, xpmx, xxfx, xxdx, xxfx, xxdx, xxd$ xpfx: /*performs transition self removal*/ $-[smxc7 \rightarrow \{AT(x, xc7, xaux1)\}xsmx]$ $-[sfxc7 \rightarrow \{AT(x, xc7, xaux1)\}xsfx]$ $-[pmxc7 \rightarrow \{AT(x, xc7, xaux1)\}xpmx]$ $-[pfxc7 \rightarrow \{AT(x, xc7, xaux1)\}xpfx]$ /*inserts an adequate transition*/ $+[smxc7 \rightarrow xsmx]$ $+[sfxc7 \rightarrow xsfx]$ $+[pmxc7 \rightarrow xpmx]$ $+[pfxc7 \rightarrow xpfx]\}$ ATK $(xn,xg) = \{$ ATKS $(xn,xg) = \{$

This simple example illustrates NL processing through adaptive formalisms. The following is a simplified derivation of the sentence "as antigas praças, parques e mansões restaurados" in our grammar. C ⇒₀ {A₁(C,C₁)}DC₁ ⇒₁ as pf C₁ ⇒₁ {A_{1c}(C₁, C, EP, EF)} EPF ⇒₂ as C ⇒₂ as {A₂(C,C₂,C₁)} AC₂ ⇒₃ as antigas pf C₂ ⇒₃ as antigas {A_{1c} (C₂, C, EP, EF) }EPF ⇒₄ as antigas C ⇒₄ as antigas S C₃ ⇒₄ as antigas praças pf C₃ ⇒₄ as antigas praças {A_{3c}(C₆,C₁,C₂,C,EP,EF)}EPF ⇒₅ as antigas praças C₃ ⇒₅ as antigas praças, C₄ ⇒₅ as antigas praças, S C₅ ⇒₅ as antigas praças, parques pm C₅ ⇒₅ as antigas praças, parques {A₄(C₆,EP,EM)}EPM ⇒₆ as antigas praças, parques pf C₅ ⇒₆ as antigas praças, parques e C₄ ⇒₆ as antigas praças, parques e mansões pf C₅ ⇒₆ as antigas praças, parques e mansões {A₄(C₆, EP, EF) }EPF ⇒₇ as antigas praças, parques e mansões C₃ ⇒₇ as antigas praças, parques e mansões {A₅(C₃,C₆, C₇,C_{8a},C_{8l})} A C₇ ⇒8 as antigas praças, parques e mansões restaurados pm C₇ ⇒₈ as antigas praças, parques e mansões restaurados C_{8l} ⇒₉ as antigas praças, parques e mansões restaurados.

Remark. " \Rightarrow_i ", $i \in N$, denotes derivation over the rules $P_L^i \cup P_D^i$. They are available after the execution of the adaptive actions.

3 Conclusion

Many forms are reported in the literature [5] for the representation of NLs and their processing by a computer. Extending the results achieved in our previous works, this paper reports a proposal for the implementation of a method for modeling, representing and handling context-dependencies in NLs by means of adaptive devices [3]. The incremental dynamic nature of our device turns it into an attractive and low-cost option with good static and dynamic time and space performance.

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